Training

ANSAwise - Transactions in Distributed Systems

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Abstract

Organizations wish to preserve their investments in IT, and databases are a case in point. They also wish to exploit new database technology if it is cost-effective to do so.

Real-world distributed systems will inevitably involve interfacing with databases of various types and ages (hierarchical, network, relational, post-relational/object-oriented). However, the database community has a rather different view of how distribution should be handled.

This module of the ANSAwise training programme compares and contrasts these different views, and outlines the key standards in the database world, and suggests how interfaces between distributed systems and databases can be architected.

This session ends with a discussion of open distributed transaction processing, including the CORBA transaction service.

[This is a variant of APM.1461, produced for CNET.]
Transactions in Distributed Systems

Transaction Manager (TM)

Application

Resource Manager (RM)

begin work
commit/rollback
work
requests
recovery

locks
log updates
prepare/commit/abort
In this session

- *Explain the Remote Data Access (RDA) technique*

- *Explain the Remote Procedure Call (RPC) technique*

- *Show how they can be used together to support templates for application partitioning*

- *Examine open standards in distributed database...*

- *... and in transaction processing*
Two worlds

- **Two approaches to distributed systems are predominant**
  - Remote Data Access (RDA), generally used for databases
  - Remote Procedure Call (RPC), generally used for distributed computing
- **RDA is based on the SQL structured query language**
  - usually to access a relational database
- **RPC is based on IDL interface definition languages**
  - usually to access an application server
- *Both are ‘client-server approaches’...*
- *...Each has unique advantages*
  - How can applications take advantage of both?
RPC and RDA for application servers

• **RPC from first to second tier**

• **RDA from second to third tier**

• *there are other possible combinations too*
Remote Data Access (RDA) Overview

- **Application interacts with database using SQL statements**
- **Statements can retrieve, update, delete, or insert records:**
  - a single record
  - a set of records at once
  - a set of records, one at a time
- **Statements can be grouped into transactions**
  - a transaction succeeds or fails as a whole
- **The location of the remote database is transparent**
RDA technique

• **Basic Features**
  - Transparent remote SQL
  - Dialog interaction
  - Dynamic SQL

• **Enhancements**
  - Distributed transaction processing
  - Distributed query processing
  - Static SQL
  - Stored procedures
RDA Advantages

- **Flexibility for clients to define application-specific views and high level operations**

- **Vendor-supported interfaces to application development tools**
  - data browsers
  - report writers
  - 4GLs
  - spread sheets
  - graphing packages
  - ...and so on
Remote Procedure Call (RPC) Overview

- Application interacts with object using a remote procedure call defined in IDL

- Statements invoke individual operations defined in IDL

- Remote procedure call can do anything an ordinary procedure call could do

- Transaction services can be used if available

- The location of the remote object is transparent
RPC technique

- **Basic features**
  - transparent remote procedures
  - “one-shot” interaction
  - static procedures

- **Enhancements**
  - “transactional RPC” (dialog interaction)
  - distributed transaction processing
  - dynamic procedure definition
RPC Advantages

• **High level of abstraction**
  - client application programmer convenience
  - semantic integrity protection
  - low communications traffic

• **Static compilation**
  - efficiency
  - automated type checking

• **Applicability to non-database services**
  - for example, device control
IDL and SQL

• **IDL defines only the interface; the logic can be written in any programming language**

• ... **SQL defines the logic; it can be executed**

• **There are various different IDL syntaxes (for example, CORBA IDL and DCE IDL)...**

• ... **there is one SQL syntax**
• A simple service (Echo)

interface Echo {

    // Comment lines start with two slashes

    string Echo (in string Src);

    string Reverse (in string Src);

};
SQL (Structured Query Language)

- **A simple SQL statement**

  ```sql
  SELECT employee_name, title, commission
  FROM employees
  WHERE employee_name IN
    (SELECT employee_name
     FROM employees
     WHERE commission > 8000
     AND title = 'account_manager');
  ```
Evolution of SQL

- **1986: SQL-86**
  - now obsolete

- **1989: SQL-89**

- **1992: SQL-92**
  - Entry SQL: roughly equivalent to SQL-89
  - Intermediate SQL: what most database vendors already supported
  - Full SQL: other well-understood features

- **1998?: SQL3**
  - parts of SQL3 will appear earlier, as addenda to SQL-92
Some features of SQL3

• **Stored procedures**
  - ‘computational completeness’
  - triggers and assertions

• **Object data management**
  - encapsulation, polymorphism, inheritance
  - multiple inheritance, generalization/specialization hierarchies
SQL Diversity

- Each vendor’s SQL has restrictions and extensions
  - you will probably have to suffer some restrictions and use some extensions...
  - ...avoiding vendor lock-in requires care
Date’s 12 Rules for Distributed Database - Autonomy

• 1. Local autonomy
  - administrators retain control over their local databases

• 2. No reliance on a central site
  - all processing is decentralized

• 3. Continuous operation
  - cannot shut down the entire system
Date’s 12 Rules for Distributed Database
- Distribution Transparency

- 4. Location independence (transparency)
- 5. Fragmentation independence (transparency)
- 6. Replication independence (transparency)
- 7. Distributed query processing (optimization)
- 8. Distributed transaction management
Date’s 12 Rules for Distributed Database - Heterogeneity

- **9. Hardware independence**
  - one distributed database can include different hardware platforms

- **10. Operating system independence**
  - one distributed database can include different operating systems

- **11. Network independence**
  - one distributed database can span LANs/WANs with different network topologies and protocols

- **12. DBMS independence**
  - one distributed database can include DBMSs from different vendors
Distributed Database

- This template can support multiple databases transparently
Distributed Database in the Real World

- Much R&D effort has been sunk into distributed databases
- No vendor has yet delivered a full distributed database solution
- ...However, practical partial solutions do exist
Distributed Database viewed by Distributed Systems

- **The 12 rules imply complete (non-selective) distribution transparency**
  - many applications don’t require this; a partial solution may be enough
  - performance may be unacceptable
  - it is unlikely to be fully achievable, anyway

- **There’s an implicit assumption that all the databases are relational**
  - this assumption would be unrealistic; 80% or more of a typical organization’s data is in non-relational form

- **There’s a need to consider integration with emerging ‘data persistence’ technologies**
  - post-relational and object databases
  - non-database distributed object stores
Database Middleware

- “Database middleware” products currently aim at a rather different set of issues from distributed systems
  
  - heterogeneous remote data access
  
  - gateways to mainframe databases
  
  - database-to-database extraction

- Products supporting query access only (not updates) already provide reasonable performance and stability
Microsoft ODBC (Open Database Connectivity)

- *De-facto specification; ODBC 1.0 specification issued in 1991*
- *Tackles the problem of heterogeneous remote data access*
- *Widespread support from database vendors*
- *Widespread support from application vendors (at least 170 products)*
- *Supports non-relational as well as relational data*
- *Also incorporated into Apple Data Access Manager (DAM)*
Distributed Systems viewed by Distributed Database

- *Any application partitioning introduces complexity*

- *RPC isn’t necessary for small-scale client/server systems*

- *RPC is useful across a broad range of applications*

- *RPC/application server development products are unproven*
  - products are immature...
  - ...vendors must provide evidence that their products enable rapid development of large-scale heterogeneous client/server systems
Transactions - the ACID properties

- **Atomicity**
  - all-or-nothing

- **Consistency**
  - transactions must be self-contained logical units of work

- **Isolation**
  - concurrent transactions must not affect each other

- **Durability**
  - what’s done must not be undone
Distributed Transaction Processing

• In a distributed system, how are transactions involving multiple databases coordinated?

• What about transactions involving operations other than database update?
Typical Open TP Core Components (much simplified)

- Application
  - begin work
  - commit/rollback
  - work

- Transaction Manager (TM)
  - requests
  - recovery
  - recovery
  - locks
  - log updates
  - prepare/commit/abort

- Resource Manager (RM)
X/Open and ISO TP Model

Transaction Manager (TM) → OSI/TP, CCR → Transaction Manager (TM)

Application

Transaction Manager (TM)

Resource Manager (RM)

Communications Resource Manager (CRM)

Resource Manager (RM)

Server

TX

XA

XA+

XATMI

TxRPC
Example Transaction Managers

- Encina (Transarc, now owned by IBM)
- Tuxedo (Unix Systems Laboratories, now owned by Novell)
Encina (IBM/Transarc)

- Application Interface - Transactional C (proprietary language extension)
- Transaction Model - Nested
- Communications - Transactional RPC (TRPC)
- Infrastructure - OSF DCE
Tuxedo (Novell/Unix Systems Laboratories)

• Application Interface - ATMI (C, C++, COBOL, 4GLs)

• Transaction Model - Flat

• Communications - Tuxedo

• Infrastructure - Most Unix
CORBA Transaction Service

- **Is an Object Service**
  - see *CORBA Services* by the Object Management Group (Wiley)
- **Specifies transactional and recoverable objects**
  - similar to Resource Managers in the X/Open DTP model
- **Tracks object usage automatically across threads**
- **Supports flat and nested transactions**
- **Supports legacy TP standards and protocols**
  - see Appendix A of the Transaction Service Specification
Summary

- **RPC and RDA can be used together to build large-scale client/server systems**

- **Distributed databases, even incomplete, may be part of the solution**

- **Distributed databases alone are unlikely to be sufficient for large-scale client/server systems**

- **Database standards are continuing to evolve rapidly**

- **Distributed transaction processing support is needed whether RPC, RDA, or distributed databases are used**