



**Poseidon House
Castle Park
Cambridge CB3 0RD
United Kingdom**

TELEPHONE:
INTERNATIONAL:
FAX:
E-MAIL:

**Cambridge (01223) 515010
+44 1223 515010
+44 1223 359779
apm@ansa.co.uk**

Training

ANSAwise - The ODP Computational Model [Eurocontrol]

Mark Madsen

Abstract

This presentation is a module of the ANSAwise training program.

It describes the following concepts in the ODP/ANSA Computational Model.

- Objects
- Encapsulation
- Shared state
- Multiple interfaces
- Service configurations
- Operations and Terminations
- Interface types
- Type conformance

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Briefing Note

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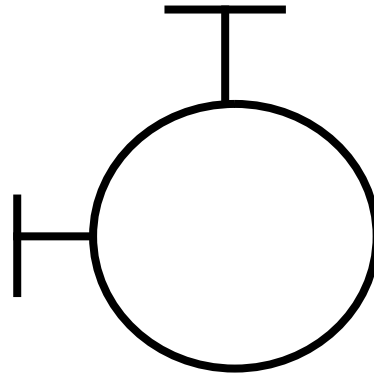
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The Computational Model





In this session

- Explain the concepts of the Computational Model
- Explain the rules that are imposed, and their importance
- Explain how objects and interfaces are used



What is the Computational Model?

- It is a model for specifying service provision and use
 - Objects may provide multiple services to other objects
 - Objects may use multiple services from other objects
- Objects and their services may be created and destroyed dynamically
- The ability to use a service may be passed from one object to another
- The Computational Model is based on interacting *distributed objects*
 - It makes no assumptions about where the objects are located

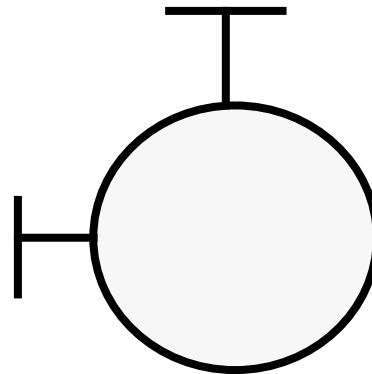


What is meant by 'object'?

- **Nowadays all software is supposed to be 'object-oriented'**
 - **OOA/OOD (analysis and design methods)**
 - **OOL (programming languages)**
 - **OODBMS (databases)**
 - **... and many more**
- **All these have one key concept in common...**

Encapsulation - the key concept

- **Objects are encapsulated**
 - every object provides a service via interfaces
 - the interface is public; the implementation is private and hidden
 - encapsulation forms a boundary; the only access to an object is via its interfaces



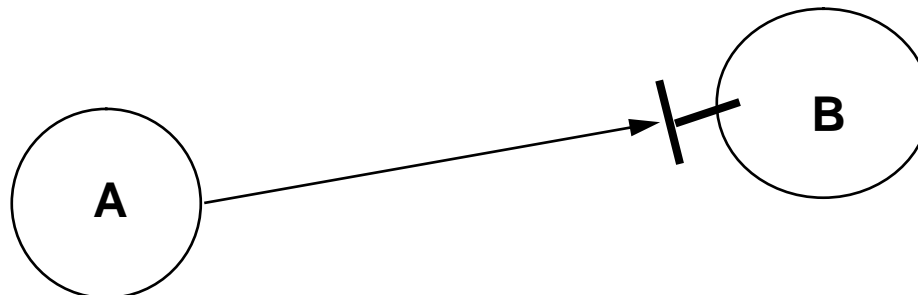


Encapsulation hides data representation

- **The boundary hides the data representation used by the object**
 - **Diversity: different objects use different representations**
 - **In a distributed system, different representations are used in different places**
 - **One object cannot manipulate another's data (only via the interface)**
 - **Data representations cannot be transferred between objects**

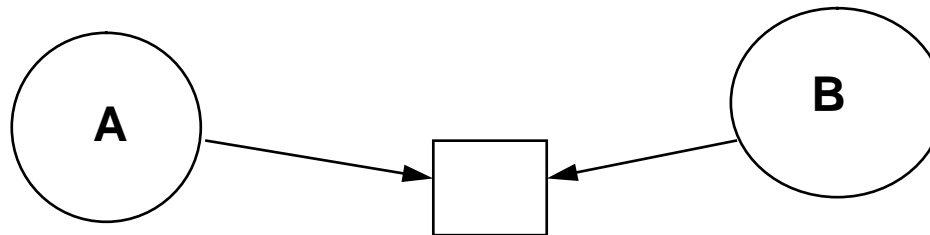
Encapsulation for distributed objects

- We cannot assume anything about the location of distributed objects
 - Objects A and B could be on the same machine, or in different countries

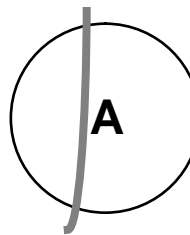


Encapsulation rules

- **Objects cannot share state directly (only access it via interfaces)**
 - *this is not allowed*

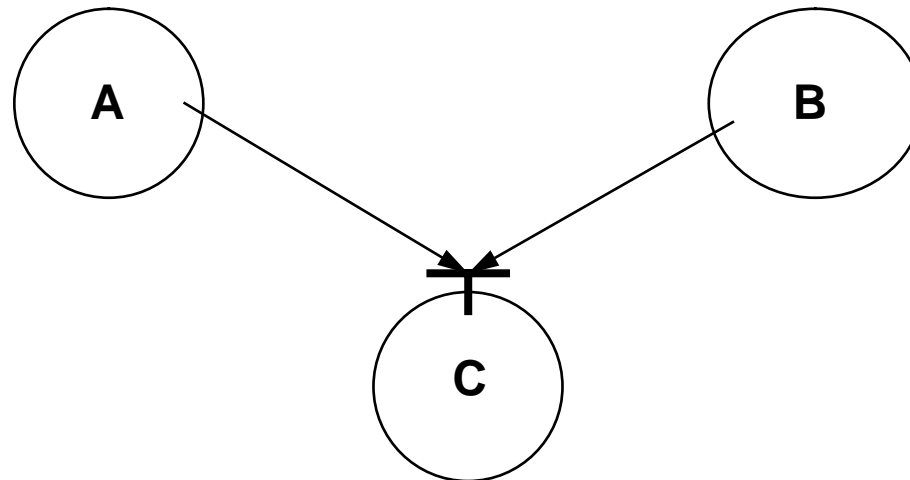


- **An object cannot be distributed in pieces; it must all be in one place**
 - *this is not allowed*



Sharing state via an object

- **No back door is needed; use an interface to a shared object instead**
 - **Both A and B can access their shared state stored in C**



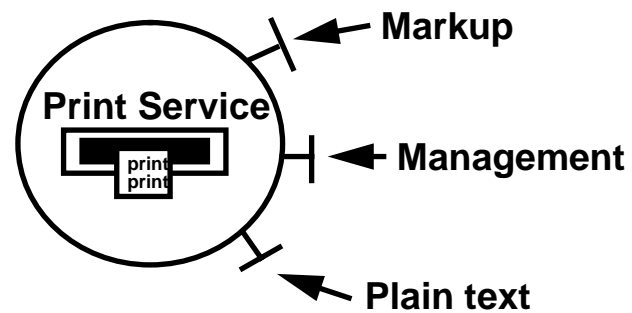


Multiple interfaces to an object

- **Multiple interfaces allow flexible configurations...**
- **... Consider a service for printing documents**
- **It must print two kinds of document**
 - **plain text documents**
 - **documents in some markup language, for example PostScript (TM)**
- **It must also have a management interface for**
 - **starting and stopping a print queue**
 - **delete print jobs**
 - **other administrative tasks**

A print service with multiple interfaces

- One object with plain text, markup, and management interfaces





Objects and Interfaces - summary so far

- **Objects encapsulate state**
 - and must take full responsibility for it
- **Interfaces provide the service access points**
- **Objects can have multiple interfaces**
- **[Object-oriented programming languages:**
 - use one construct for both encapsulation and service provision
 - provide only a single interface]

with their shared state



Operations - the actions in an interface

- An interface defines one or more *operations*
- Operations are the actions that the user of the interface can invoke
 - Similar to "methods" in object-oriented languages
- Because of encapsulation, the defined operations are the only way to act on the object
 - No “back door” [not even for management or debugging]

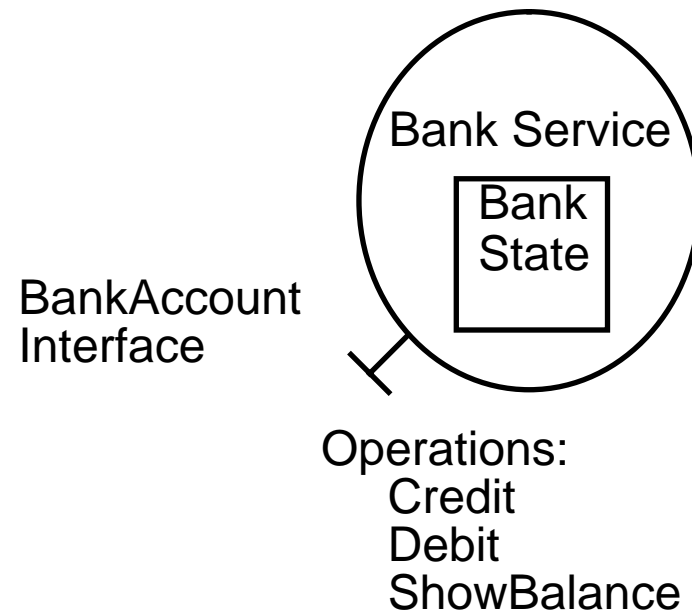


Operations - a Simple Bank example

- **Consider a bank account service...**
- **Each bank account has state**
 - **account balance**
 - **name of customer**
 - **... and others**
- **Each bank account has actions that affect this state**
 - **Credit: depositing money, increases the balance**
 - **Debit: writing a cheque, decreases the balance**
 - **Show: enquires the account balance**
- **These actions are the only way of manipulating the bank account; these actions are the *operations***



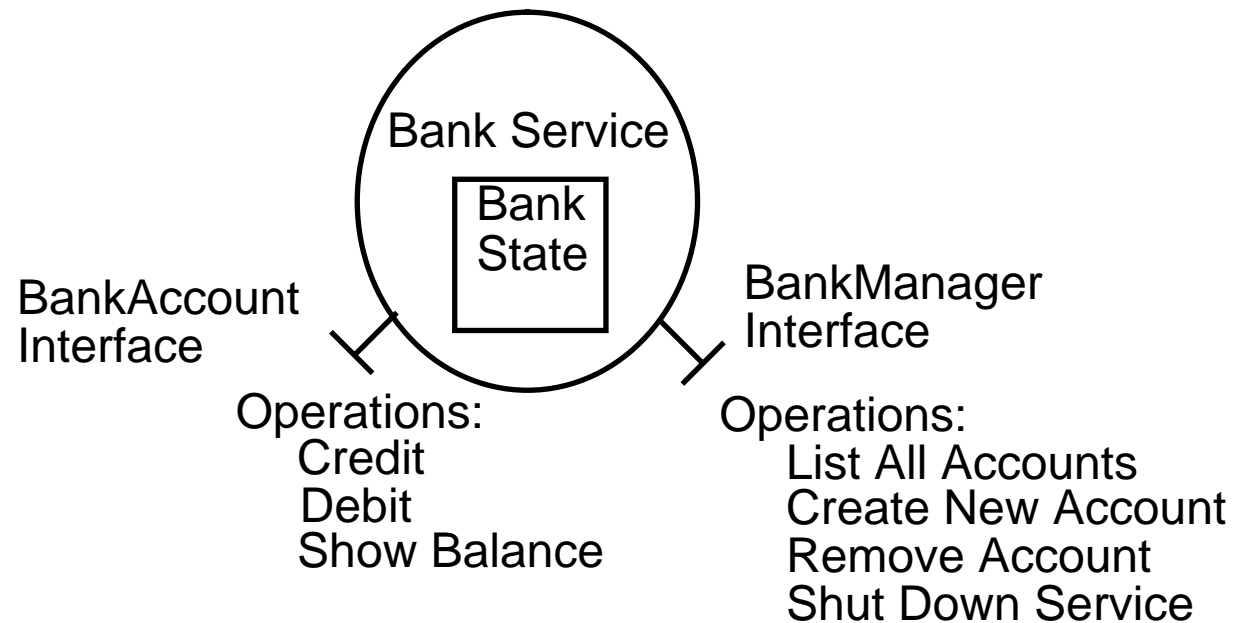
Simple Bank - Service, interface, and operations





Simple Bank - Operations in multiple interfaces

- A service can have more than one interface
- Each interface has its own operations





Operations and Arguments

- **Each operation has a name, input arguments, and output arguments**
 - **for example, the Credit operation returns the new balance**
Credit (amount : Integer) -> (newBalance : Integer)
 - **amount is a input argument; newBalance is an output argument (result)**
- **Each operation can have multiple input and output arguments**
 - **this example only has one of each**



Operations and Terminations

- An operation can have more than one possible outcome
- Consider operations on the BankAccount:
 - Credit (amount : Integer) -> (newBalance : Integer)
 - Debit (amount : Integer) -> (newBalance : Integer)
 - ... but the balance may be insufficient! - so
 - Debit (amount : Integer) -> (newBalance : Integer) ->InsufficientFunds()
- There is one Debit operation with two possible outcomes; two separate *terminations*



Named Terminations

- **The InsufficientFunds outcome is called a *Named Termination***
 - **the normal outcome is an unnamed (anonymous) termination**
 - **one termination of an operation may be anonymous**
- **Each termination may return multiple results**
- **The invoker of an operation distinguishes between the different terminations by their names**

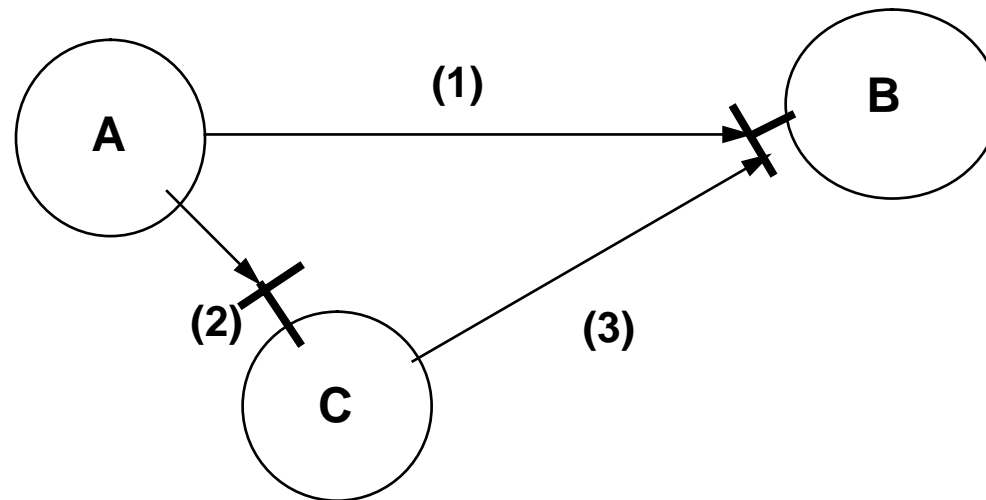


Operations, Arguments, and Terminations - summary

- **An *operation* is the unit of interaction. It has:**
 - a name
 - a signature (list of input arguments)
 - a set of *terminations* (possible outcomes)
- **A termination has:**
 - a name
 - a signature (list of output arguments)
- **Terminations make operations more powerful than ‘functions’ or ‘methods’**

Using Interfaces

- Object A is using one of object B's interfaces
- Suppose it needs to tell C to use the same interface



- It must be possible to pass a binding to B's interface between A and C

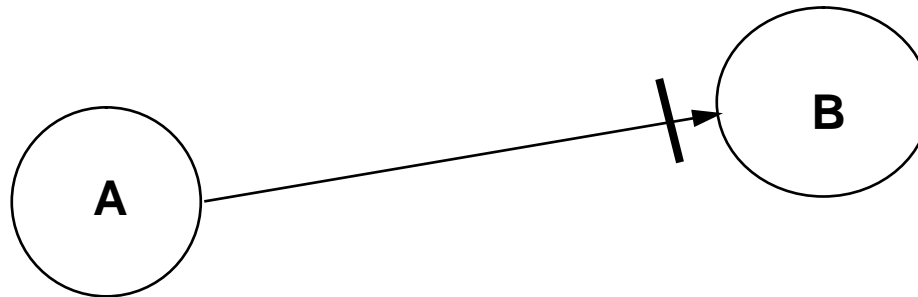


Interface Bindings

- Interfaces are identified by *interface bindings*; these can be passed to other objects as arguments (or results)
 - for example, when A passes B's interface binding to C, it calls...
C_op (b: TypeOfInterfaceToB) ->()
 - ... with B's interface binding as the argument
- Only the binding is passed, not the interface...
- ... In fact, all operations are invoked via interface bindings

The right interface?

- Suppose object A has a binding for an interface of B...
- ... how does object A know that B has the right kind of interface?
 - it can't inspect the object itself
 - objects A and B could be on the same machine, or in different countries



- It can tell, because interfaces have an *interface type*

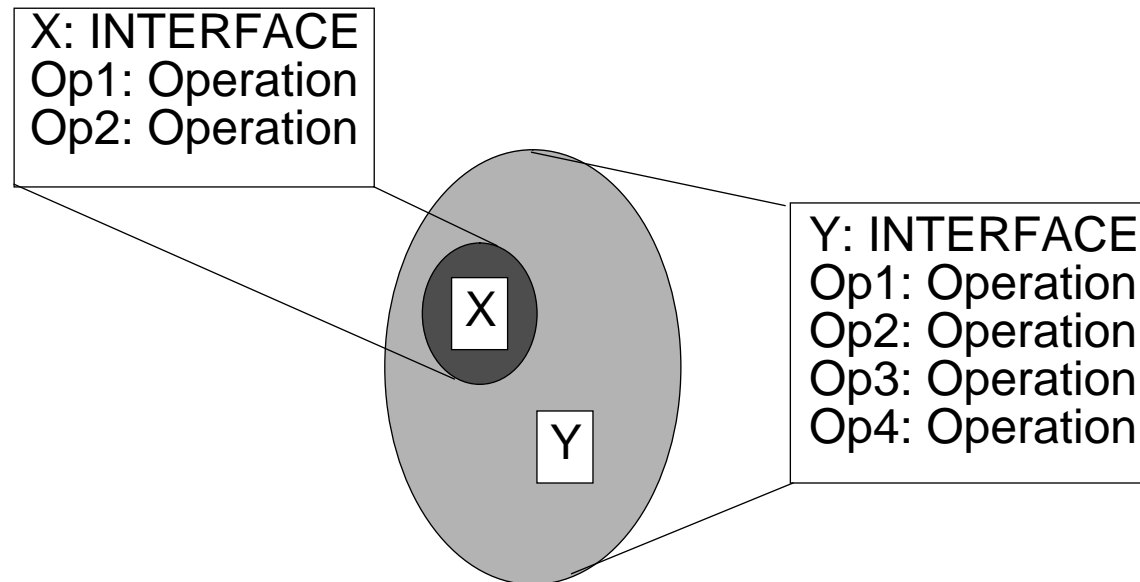


Interface types

- **The interface type consists of the set of its operations**
 - **Plus all their signatures and terminations**
- **Object A can use an interface of object B, provided that the interface provided by B *conforms* with that expected by A**
- **This means that object A must state what interface type it expects**

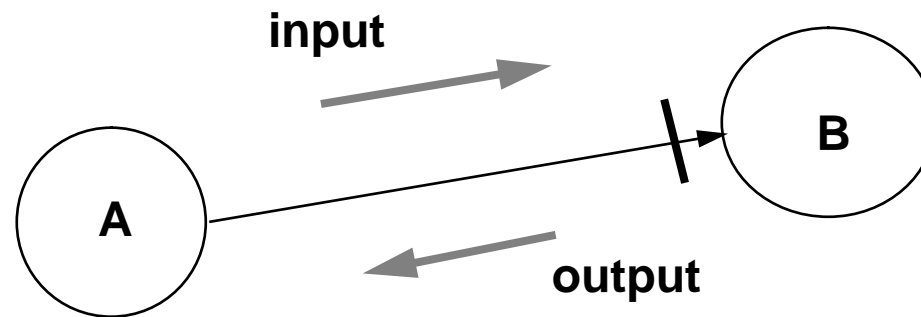
Type Conformance

- **Conforming interface types do not have to be identical...**
- **...Interface type *Y conforms to* interface type *X***



Type Conformance Has Two Sides

- To conform, B must provide at least the operations expected by A...



- ...but also, A must handle at least the terminations given by B
 - same idea, but the 'at least' rule is the other way round for terminations
- To conform, both of these must be checked... [a two-sided *contract*]
 - B must not receive unknown operations
 - A must not receive unknown terminations



Summary

- **Objects enforce strict encapsulation**
 - state cannot be shared between objects
 - only whole objects can be re-configured (*part of an object cannot be*)
- **Interfaces provide the points of service provision**
 - objects may have multiple interfaces
 - interfaces are typed
- **Interfaces are composed of operations**
 - operations can have multiple terminations with multiple results
- **For more information:**
 - for object-oriented concepts, see *Distributing Objects* (TR.018.01)
 - for a formal description, see *The ANSA Computational Model* (AR.001.01)



Computational Model - topics not covered

- **Concurrency**
 - threads and sub-threads, enforcing tree-structured concurrency
- **Streams - continuous (possibly bi-directional) information flows, for example:**
 - speech, in telephony and voice processing: an audio stream
 - video, in multimedia: a video stream
 - sensor data, in telemetry: a data stream
- **Synchronous programming constructs**
 - reactive systems, bounded execution paths
 - testing, receiving, waiting for and transmitting signals
 - watchdogs