

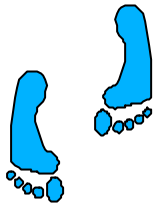
Secure Sessions

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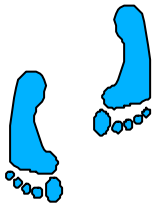
Progression



Kamae (MoD)

Secure RMI (Java)

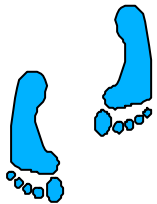
cipher and (2 level) digest sessions
active labels



E2S (Esprit)

retrofitted (in Java) onto OrbixWeb

with (2 level) digesting at application level
then translated into C++ on Orbix

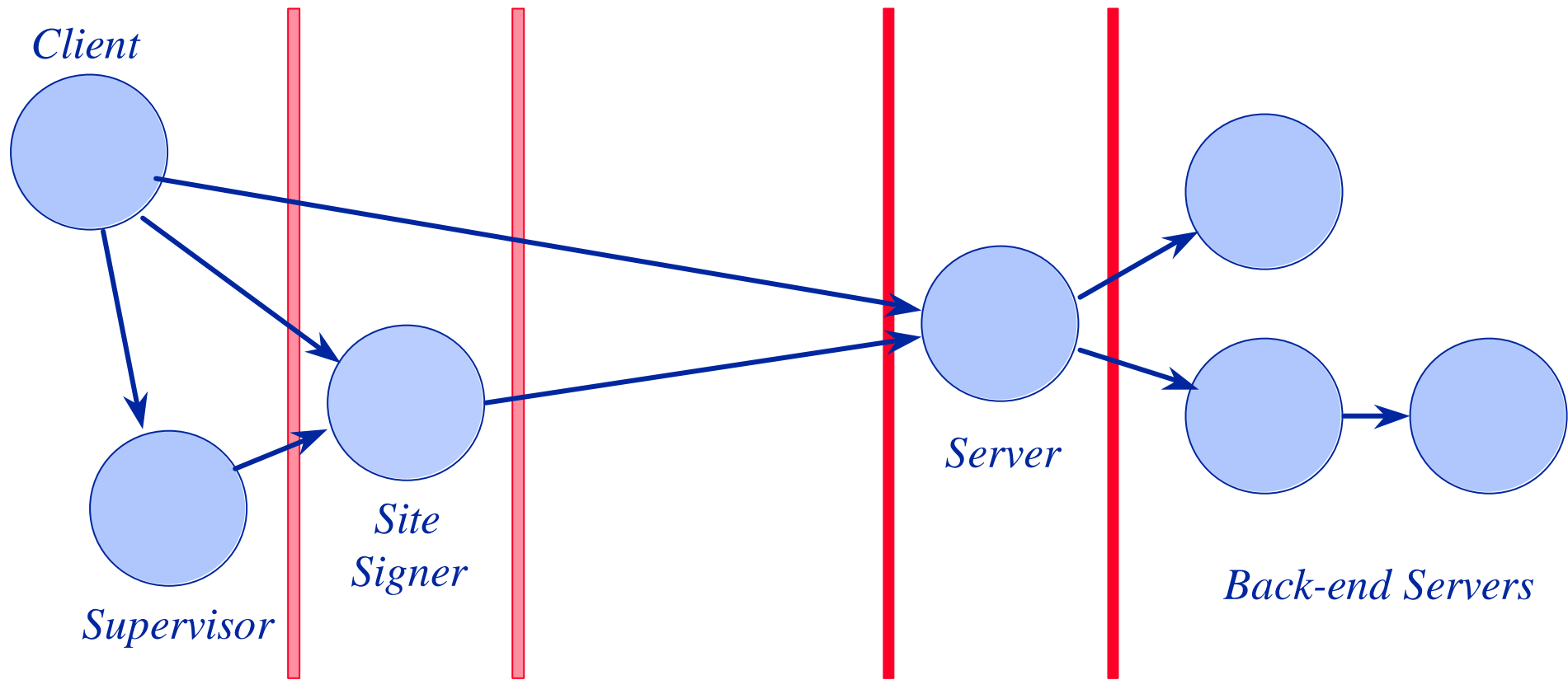


Flexinet (ANSA)

fully recursive application level digests
automatically serialized or explicitly marshalled
integrated with Flexinet cipher sessions



Securing Distributed Objects



Assumptions and Principles

- Assume necessary and sufficient infrastructure for:
 - strong object and process encapsulation
 - process encapsulation is OS dependent
 - object encapsulation is language dependent
 - application processes being responsible for their own protection
- Assume emerging smart card technology
- Maximize transparency of mechanism
- Apply same principles and protocols everywhere
 - including security service objects



Secure RMI Objectives

- Applications in total control
 - Fully transparent engineering
 - conflict handled by separating concerns
 - Allow interworking with non-secured applications
 - Allow local or remote control of policy
- in conflict

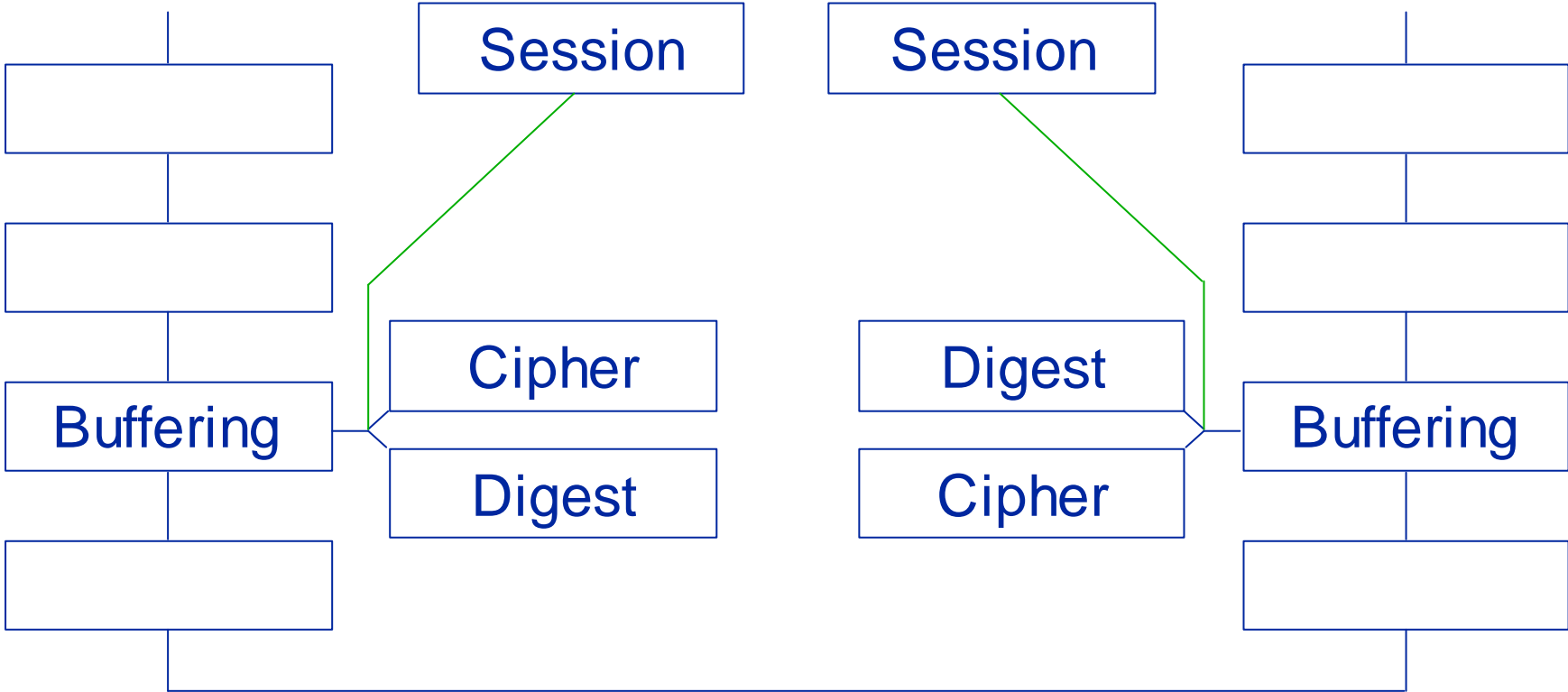


Minimal Protocol Mechanisms

- Mechanisms necessary and sufficient for purpose
 - keyed digests for integrity and authentication
 - XORed keyed digests for key distribution
 - symmetric encryption for confidentiality
 - (asymmetric signed digests for non-repudiation)
- Hash functions
 - fixed length output
 - faster
 - less legal problems in commerce (COTS)



Secure RMI Structure

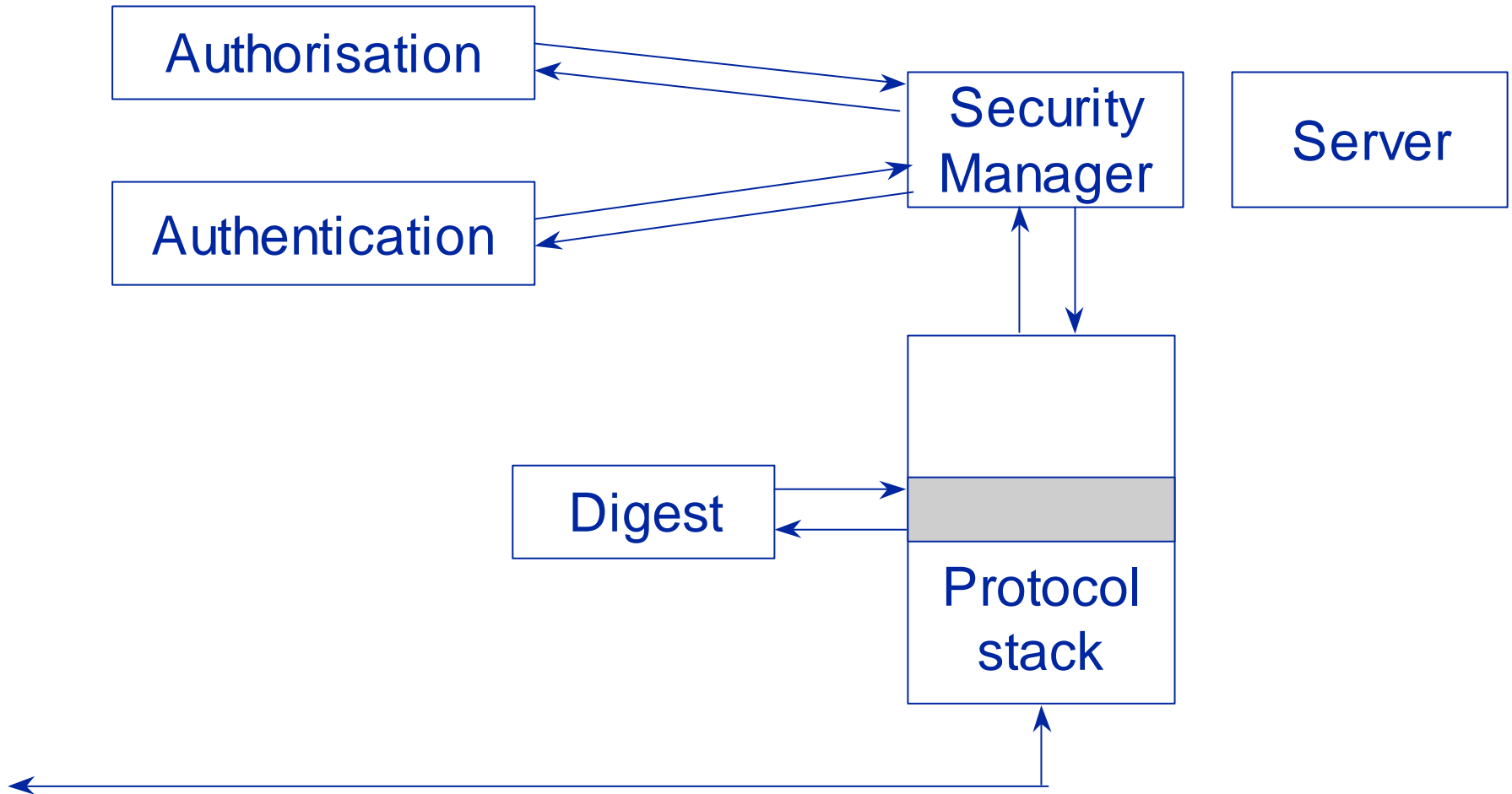


Secure RMI Engineering

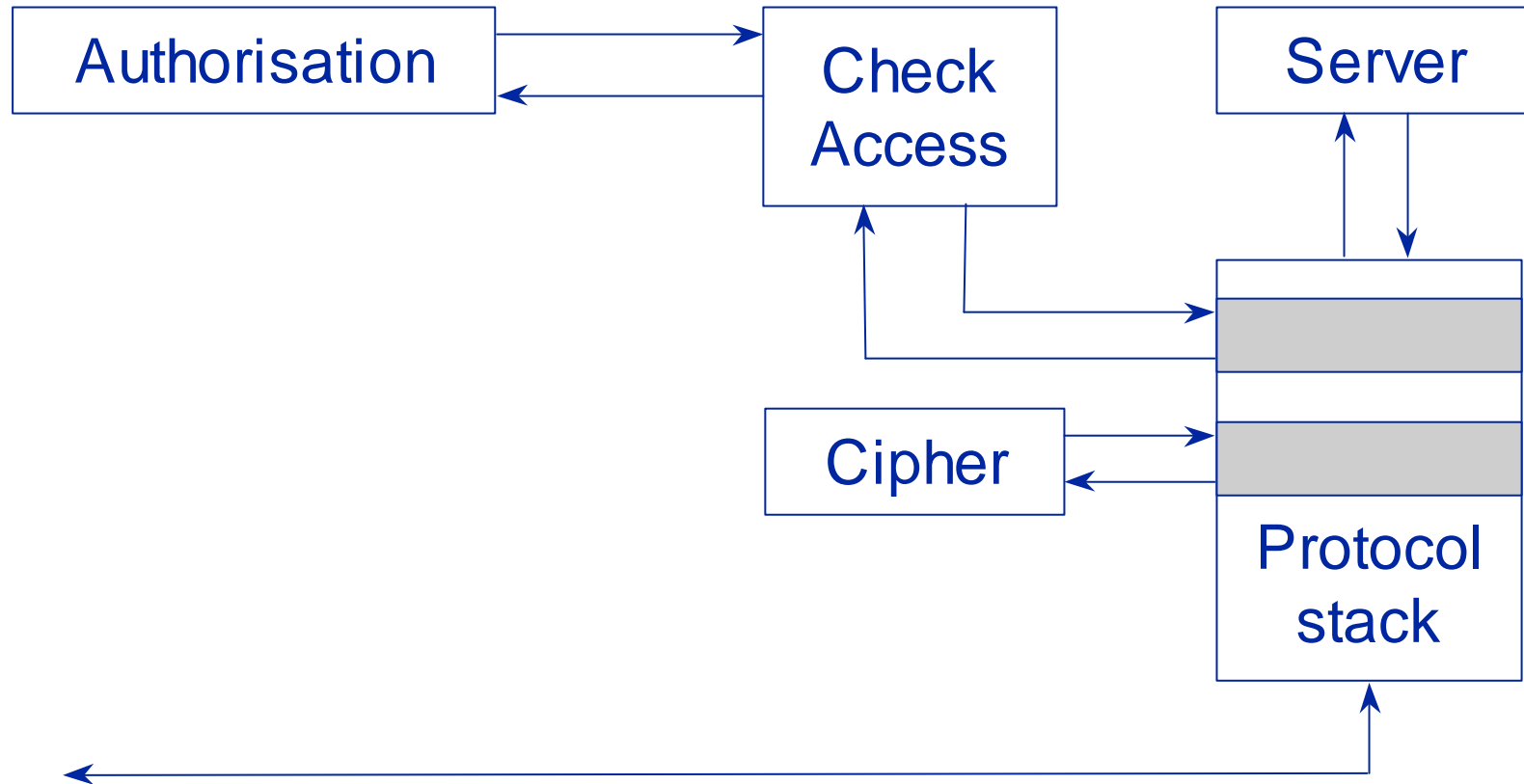
- An application security manager added to the server
- Client bound to security manager instead of server
 - use digest session to manager to open server
 - use cipher session to access server
- Server application can take control
 - or delegate to third party security servers
- No security parameters in server methods
 - check access called before each method



Opening a Server



Per Method Authorisation



Protocol Properties

- Generalisation to client-server chains
 - recursive protocol
 - session keys returned to each pair of participants
- Overlay of authentication and delegation
- Symmetric call and return engineering
- Simplifies federation of administrative domains
 - its just another indirection



Protocol Description

a -> b : $Da = [a, b, Na, --]Ka$

b -> c : $Db = [b, c, Nb, Da]Kb$

c -> k : $[c, k, Nc, Db]Kc$

k -> c : $[c, b, Nc, Nb, Kcb \oplus \langle c, b, Nc, Nb \rangle Kc, Eb]Kcb$

c -> b : $Eb = [b, c, Nb, Nc, Kbc \oplus \langle b, c, Nb, Nc \rangle Kb,$
 $[b, a, Nb, Na, Kba \oplus \langle b, a, Nb, Na \rangle Kb, Ea]Kba]Kbc$

b -> a : $Ea = [a, b, Na, Nb, Kab \oplus \langle a, b, Na, Nb \rangle Ka, --]Kab$



Lessons Learned from Kamae

- it is feasible to transparently security a remote interface
- but its hard to nest and unnest digests as byte arrays
 - and its very hard to add and extract applications arguments
- it is possible to almost hide the distinction
 - between digest and cipher references
 - but not between them and clear references
- the application managers have to be aware of digest and cipher sessions



Conclusion

Do digesting and session management at the
application level

(in management classes alongside the client and server
classes)



Migration to OrbixWeb

OrbixWeb has 3 protocol hooks
(But they don't all work with IIOP)

- Transformers
 - only access to byte stream (encryption)
 - only other info is remote host name and direction
- Per-Object Filters
 - only on server side (per method authorization ?)
- Per-Process Filters (digesting)
 - 8 versions
 - pre or post marshalled
 - client or server
 - request or reply



Approach

- encrypt in the transformers
 - cipher sessions indexed by remote host name
 - if session found then en/de-cipher
else pass in clear (for authentication digests)
 - add random IV to each cipher message
- authenticate, authorize and distribute session keys in application security managers
 - explicit marshalling of digests into byte arrays
 - explicit session setup



Invocation Path

client outRequestPremarshal

client transformer outgoing

server transformer incoming

server inRequestPreMarshal

server application

server outReplyPreMarshal

server transformer outgoing

client transformer incoming

client inReplyPreMarshal



Problems

- 1) an attacker can invoke a server in clear
- 2) if a client's host machines are behind an IP address translating firewall
 - a second client's transformer would transmit a digest request in clear
 - but server's transformer would decipher it
- 3) server needs to call back-end servers in clear



Lessons Learned from E2S

- implement sessions for all remote calls
 - add session IDs to all messages
 - plaintext, ciphertext and digested
 - can then have different session key for each user
- provide digest sessions as well as cipher sessions
 - for ensuring integrity when encryption is not allowed
 - or cooperation with the encryption mechanism is difficult
- provide flexible application level digests which:
 - can be automatically serialized or explicitly marshalled
 - can be nested to any depth
 - can enclose application data in both directions



Flexinet

- sessions are fully engineered into the protocol stack
- CipherSession & DigestSession are subclasses of Session
 - no other security engineering in the protocol stack
- fully nested application level protocol for authentication, authorization and key distribution using keyed digests
 - using classes which can be serialized or explicitly marshalled
 - and can carry application data in both directions
- explicit binding of CipherSession & DigestSession
 - interfaces bound to a CipherSession or DigestSession cannot be bound to any other type of session

